FlexPS: A flexible and scalable Parameter Server for Distributed ML

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https://github.com/Yuzhen11/flexps

4 Needs for ML & AI Tech

Resource Efficient: Use less CPU, GPU, memory, networks ...

Scalable: Linear performance increase with more computation resource

Correct: Can we trust the result?

General ML/AI Platform: Deep learning is only 10%

2



Big Data

facebook

1B+ Users 30+ Petabytes



32M+ pages

WIKIPEDIA The Free Encyclopedia



100h+ video uploaded every minute



645 millions users 500 millions tweets / day

Big Model

Large Model is better for Big Data

Google Brain deep learning for images:

1 ~ 10 billion model parameters

Netflix collaborative filtering for video recommendation:

1 ~ 10 billion model parameters

Topic Models for text analysis:

Up to 1 Trillion model parameters

stem

auxiliary classifiers

Movie Ratings	Zora	Sophie	Jordan	Ernie	Christie
Harold and Kumar Escape	8	4	?	?	4
Ted	?	?	8	10	4
Straight Outta Compton	8	10	?	?	6

And the second s

Iterative Convergent Algorithm

Data: D

Model *L* (ie. a fitness function such as likelihood)

Algorithm: update the model's parameters **A** iteratively until it converges

$$A^t = F(A^{t-1}, \Delta_L(A^{t-1}, D))$$

Algorithms: Stochastic Gradient Descent (SGD), Stochastic Average Gradient (SAG), stochastic variance reduced gradient (SVRG).

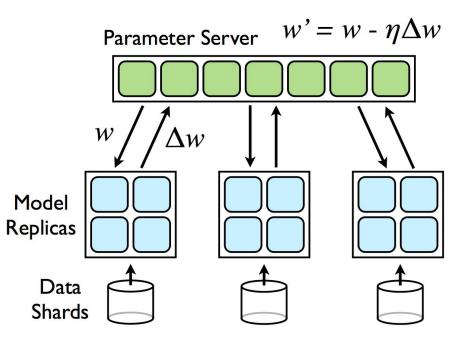
Models: Logistic Regression, Support Vector Machine, Kmeans, Neural Network

Parameter Server

separate the working units into *workers* and *servers*

- Parallel workers update models stored distributedly in servers

- Easy-to-use key/value store interface

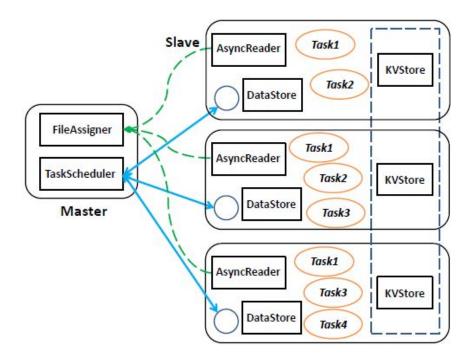


FlexPS: Architecture

FlexPS is organized in a master/slave architecture

- Master
 - TaskScheduler
 - FileAssigner

- Slave
 - KVStore
 - DataStore
 - AsynReader
 - Tasks



Architecture: TaskScheduler

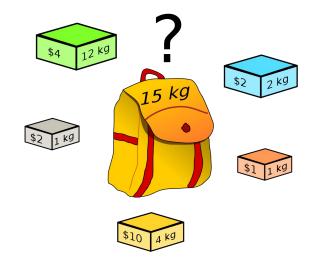
TaskScheduler decides which task to run & where it runs.

3 scheduling algorithms:

- Sequential scheduling
- Greedy scheduling
- Prioritized scheduling

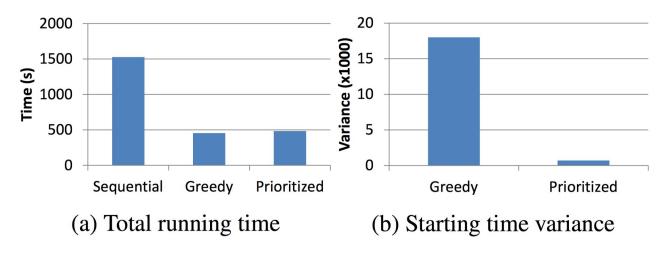
Optimal scheduling algorithm?

Multiple Knapsack Problem \rightarrow NP Hard



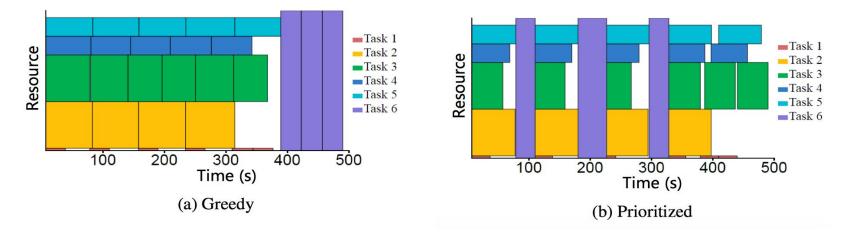
Evaluation: Task Scheduling

- **6** tasks with different workloads, **300** threads
- Greedy and Prioritized significantly outperform the Sequential due to task parallelism
- Prioritized: prevent starvation & high throughput



Evaluation: Task Scheduling

- 6 tasks with different workloads, 300 threads
- Resource \rightarrow # of threads
- Both algorithms achieve high resource utilization at most time
- Prioritized scheduling eliminates starvation (Task 6)



Architecture: KV-Store

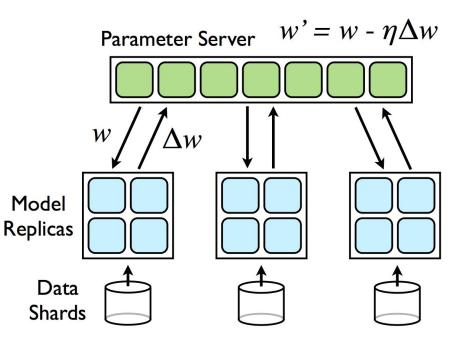
Programming Model:

Get(keys)

Add(keys, vals)

GetChunk(keys)

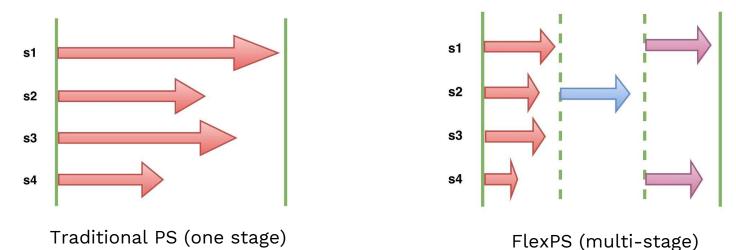
AddChunk(keys, vals)



Key Design: Multi-stage ML

Breaks down a machine learning task into different stages

A *stage* runs a sub-task on a specific set of computing resources characterized by the number of slaves and the location of these slaves



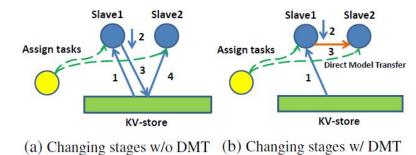
Optimization: Direct Model Transfer

Normal Procedure:

- 1. Load the model from KV-Store
- 2. Perform training locally
- 3. Dump the model to KV-Store
- 4. Repeat above steps in subsequent stages

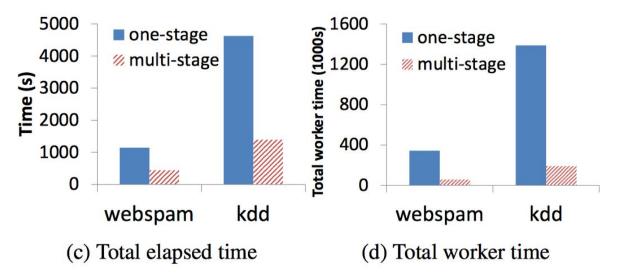
Direct Model Transfer:

- Bypass KV-Store
- Boost the performance of changing stages by 23%



Evaluation: LR Performance on SVRG

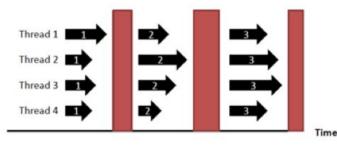
- Total worker time: the total amount of (worker × time) in all stages (reflect total CPU hours)
- 3.5x on elapsed time, 8x on worker time



Background: Consistency models

Bulk Synchronous Parallel(BSP)

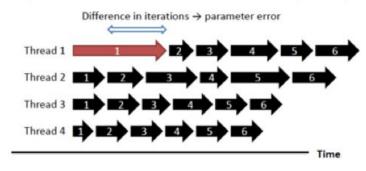
- Synchronization Barrier (Parameters read/updated here)
- (a) Machines perform unequally
- (b) Algorithmic workload imbalanced



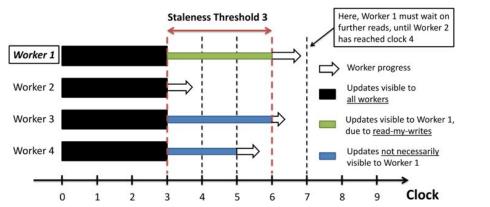
Asynchronous

Parameters read/updated at any time

Asynchronous is fast but has weak convergence guarantees



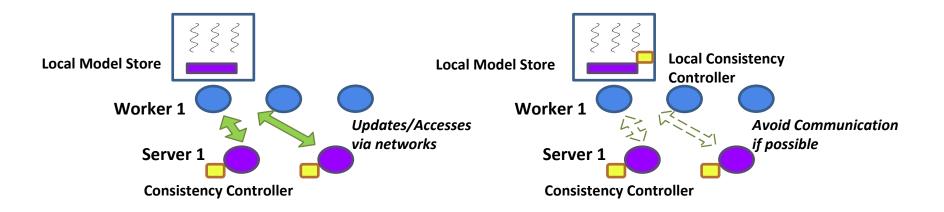
SSP: Bounded Staleness and Clocks



Ho, Q. (NIPS 2013). More E ective Distributed ML via a Stale Synchronous Parallel Parameter Server.

Key Design: Local Consistency Control

Pop up consistency controller from server side



Key Design: Local Consistency Control

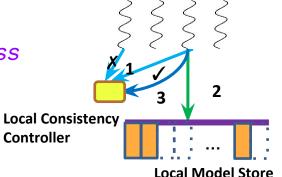
Model access procedure:

- 1. working thread checks with the local consistency controller
- 2. If the consistency requirement is satisfied, this thread

is granted to access the model

3. When slowest thread finish, update *min_progress*

& notify blocked threads



All updates & consistency controls are performed locally!

Key Design: Local Consistency Control

Concurrency control:

- Multiple threads may access the shared local model simultaneously
- Sort & divide the local model into chunks
- Attach a lock (mutex) to each chunk
- Tradeoff: Lock every single parameter vs Lock a chunk
- Experiments: Less contention, fast model access (1-2% of total time)

FlexPS: Optimizations

Load Balancing:

A global *FileAssigner* to keep the block information of all files in HDFS Assigns blocks to the worker threads according to *data locality*

Data loading on-the-fly:

AsyncReader module with the classical producer-consumer paradigm

reader threads load the data from HDFS and store them to a pre-allocated buffer

worker threads consume the data in the buffer and use the data to train the model

Comparison: LR with SVRG

FlexPS-Opt: offline search for optimal stage config

FlexPS-Auto: runtime search for stage config

FlexPS-: disable the flexible parallelism control of FlexPS

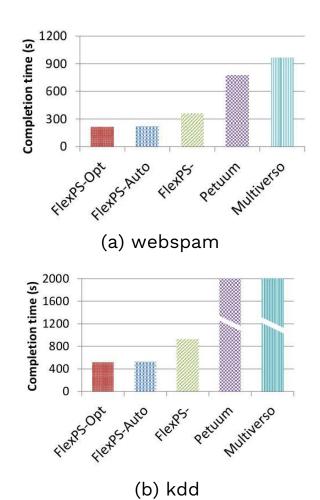
Dataset: webspam, kdd

Batch size of the stochastic step:

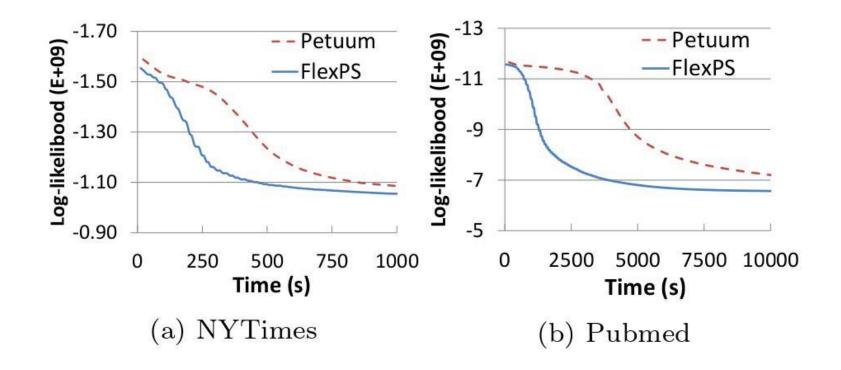
0.1% for webspam and 0.001% for kdd

Number of training epochs: 10

2-5 times speedup in task completion time!



Comparison: Convergence speed on LDA



Conclusion

We proposed FlexPS, which provides a more *flexible* and *scalable* PS framework for distributed ML with innovative designs:

- TaskScheduler: flexible task scheduling (sequential, greedy, prioritized)
- Multi-stage design: make the most of the computing resources
 - Direct model transfer: Bypass the KV-Store
- Local consistency control: Avoid network traffic
 - Concurrency control: Lock single parameter vs Lock a chunk
- Load Balancing: exploit the data locality
- Data load on-the-fly: asynchronous data loading



Example: Logistic Regression

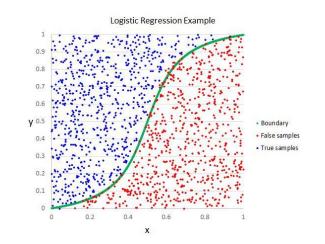
y = w.X + b

y: Discrete class [0, 1] prediction

X: data set (a matrix)

Cost: Correct/Wrong prediction from actual

Goal: Find scalars w, b



$$egin{aligned} P(y=1|x) &= h_ heta(x) = rac{1}{1+\exp(- heta^ op x)} \equiv \sigma(heta^ op x), \ P(y=0|x) &= 1-P(y=1|x) = 1-h_ heta(x). \end{aligned}$$

Iterative Convergent Algorithm, often solved by SGD

Example: Stochastic Variance Reduced Gradient

Procedure SVRG

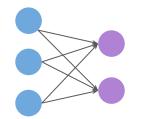
```
Parameters update frequency m and learning rate \eta
Initialize \tilde{w}_0
Iterate: for s = 1, 2, \ldots
  \tilde{w} = \tilde{w}_{s-1}
  \tilde{\mu} = \frac{1}{n} \sum_{i=1}^{n} \nabla \psi_i(\tilde{w})
  w_0 = \tilde{w}
  Iterate: for t = 1, 2, ..., m
     Randomly pick i_t \in \{1, \ldots, n\} and update weight
       w_t = w_{t-1} - \eta(\nabla \psi_{i_*}(w_{t-1}) - \nabla \psi_{i_*}(\tilde{w}) + \tilde{\mu})
  end
  option I: set \tilde{w}_s = w_m
  option II: set \tilde{w}_s = w_t for randomly chosen t \in \{0, \ldots, m-1\}
end
```

Example: Multi-stage LR with SVRG

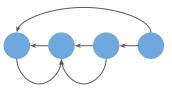
```
/* Step 1: define the stages */
auto fgd lambda = [](Info info) {
// Get model parameter w from the KV-store
// Calculate gradient
// Update full gradient u in the KV-store
};
auto sqd lambda = [](Info info) {
// Get w and u from the KV-store
// Calculate gradient
// Perform variance-reduced update
// Update w to the KV-store
};
/* Step 2: set the parallelism degree */
MultiStageTask task;
task.SetStages({{fgd lambda, 100}, {sgd lambda, 10}});
/* Step 3: submit the task */
engine.SubmitAndWait(task);
```

Machine Learning Systems

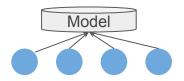
Iterative MapReduce



Data Flow Graph



Parameter Server











dmlc **mxnet**

Why Parameter Server?

Support for big model

The driver can't hold big model due to limited memory size

Flexible Consistency Control

Allows asynchronous operation

Spark and Hadoop are Bulk Synchronous Parallel

Better network utilization, and lets you scale your models

Existing PS Systems

ps-lite:

Underlying communication module of *MXNet*

Flexible consistency models, Fault Tolerance

Petuum:

Developed by SAILING Lab from CMU, now a start-up company

Flexible consistency models (SSP), Model parallelism

Multiverso:

Core module of the Microsoft Distributed Machine Learning Toolkit (DMTK) Abstract Communication APIs, Supports for deep learning systems (torch, theano)

Limitations in Existing PS Systems

Flexibility in Parallelism Control

Machine learning algorithms may have varying workloads in different training stages

Stochastic algorithm: smaller mini-batch sizes in earlier stages for fast convergence and larger mini-batches later to avoid oscillation

variance reduction algorithms: repeats the following two phases until convergence:

- 1. Compute the full gradient using all data records
- 2. Update the parameters in a stochastic manner

Existing PS Systems only allow a *fixed* number of worker threads (i.e., fixed parallelism) throughout the whole training process

Limitations in Existing PS Systems

Optimization for Sparse Data

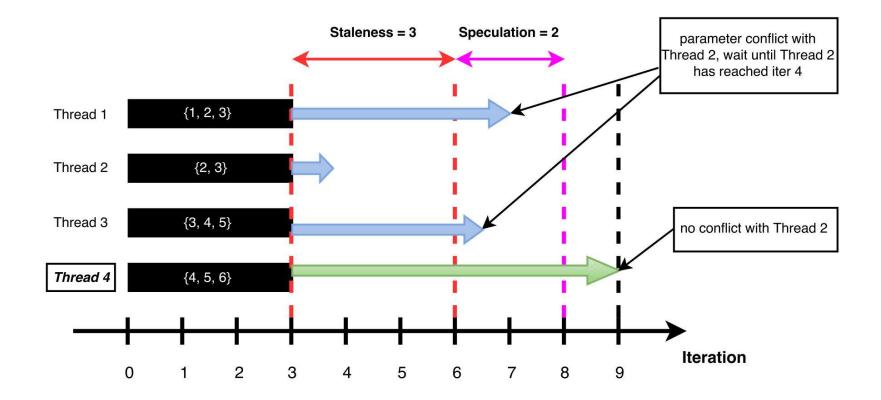
In most cases of the high dimension dataset, the data is actually sparse (data missing).

Only store the non-zero entries as key-value pairs to save memory.

Special optimization for sparse machine learning problem with high-dimensional parameters needs to be proposed for lower computational/communication cost.

Existing PS Systems typical support communication models like BSP/SSP/ASP, which make no difference between sparse and dense data.

FlexPS: SparseSSP



FlexPS: SparseSSP

Further optimize SSP for sparse data

Allow worker move to next iteration if the set of parameter that it is manipulating does not conflict with the slowest worker's

Max clock difference between workers:

Staleness + Speculation

Algorithm 1 SparseSSP Algorithm to check if worker i can move on to next iteration

Ree	quire: w_i : worker i c_i : the clock value w_i s: staleness sp: specu-
	lation P_i : the set of parameter that w_i is manipulating
1:	initialize flag = True
2:	for every worker w_i from all worker where $i \neq j$ do
3:	if $c_i - c_j \leq s$ then
4:	continue;
5:	else if $s < c_i - c_j <= s + sp$ then
6:	if $P_i \cap P_j = \emptyset$ then
7:	continue;
8:	end if
9:	end if
10:	flag = False;
11:	break;
12:	end for
13:	if flag = True then
14:	w_i can move on
15:	else
16:	w_i waits until the slowest worker finish its current iteration
	1.0

How often will there be conflict?

Total number of parameter of dataset: n

Number of non-zero parameter in each training batch: m

Sparsity: m/n

Probability of parameter conflict between two workers: $P(conflict) = 1 - \frac{\binom{n-m}{m}}{\binom{n}{m}}$

Dataset	kdd	avazu	criteo	url
# of features	5.5x10^7	10^6	10^6	3.2x10^6
# of samples	1.5x10^8	4.0x10^7	4.6x10^7	2.4x10^6
Sparsity	2.0x10^-7	1.5x10^-5	3.9x10^-5	3.1x10^-5
Pr of conflicts	2.2x10^-6	2.2x10^-4	1.5x10^-3	3.0x10^-3

Evaluation: SparseSSP

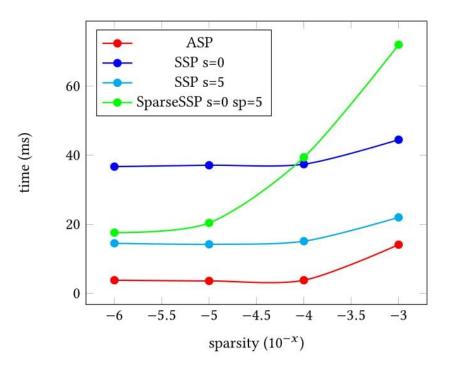


 Table 2: Experiment parameter

# of non-zero	10	100	1000	10000
Sparsity	10^{-6}	10^{-5}	10^{-4}	110^{-3}
Pr of conflicts	10^{-5}	10 ⁻⁴	$9.5x10^{-2}$	0.9955

Figure 10: Comparison of different parallelism protocol in sparse data

Future Work: Task based PS

Straggler Problem

The slowest worker in a cluster, cripples the scalibility of ditributed ML

Straggler mitigation via tiny tasks:



High level & functional: (x;w) => delta

Abstract the threads, user only focus on defining ML task

No need for SSP: tiny tasks --> load balancing

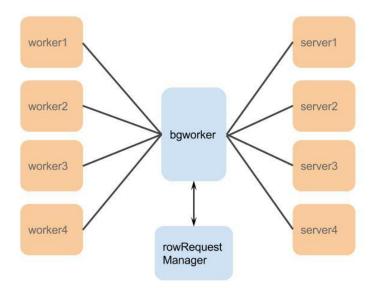


Process Cache

Currently, the worker need to fetch parameter from the server directly.

A cache could be created in the worker node to cache some temporary parameter.

It accelerate getChunk or AddChunk operations.



Load Balancing

Currently, load balancing is done during the loading data task, you could manually allocate data to different worker nodes.

In the future, load balancing could be done during the execution of task:

- 1. The worker threads will steal some data from other workers when it is blocked by consistency control.
- 2. The fast worker will help slow worker once the fast worker finish its own task.

Backup slides

